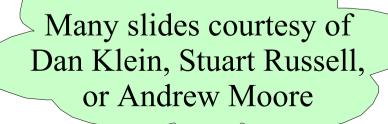
# **Adversarial Search**

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- CS 421: Introduction to Artificial Intelligence
- 9 Feb 2012





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CS421: Intro to AI

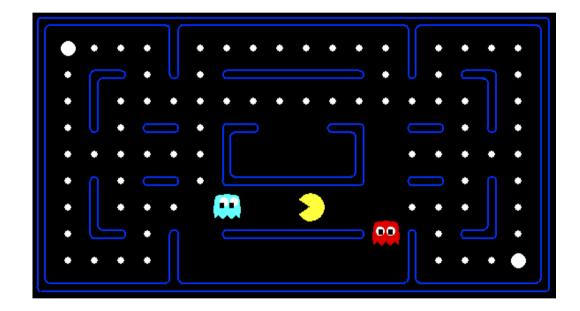
### Announcements

#### None

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### **Adversarial Search**



[DEMO: mystery pacman]

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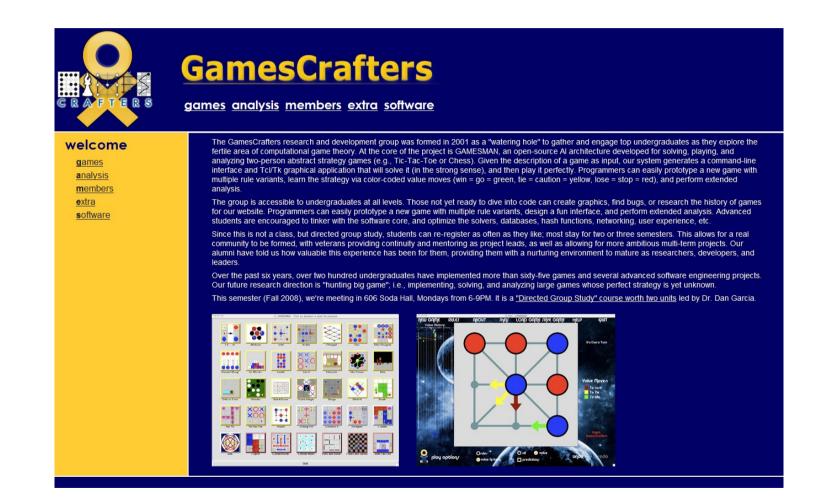
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# Game Playing State-of-the-Art

- Checkers: Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994. Used an endgame database defining perfect play for all positions involving 8 or fewer pieces on the board, a total of 443,748,401,247 positions. Checkers is now solved!
- Chess: Deep Blue defeated human world champion Gary Kasparov in a six-game match in 1997. Deep Blue examined 200 million positions per second, used very sophisticated evaluation and undisclosed methods for extending some lines of search up to 40 ply.
- Othello: human champions refuse to compete against computers, which are too good.
- Go: human champions refuse to compete against computers, which are too bad. In go, b > 300, so most programs use pattern knowledge bases to suggest plausible moves.
- Pacman: unknown

### GamesCrafters

#### http://gamescrafters.berkeley.edu/



# **Game Playing**

Many different kinds of games!

**Examples?** 

- Axes:
  - Deterministic, 1 player, perfect information?
    Deterministic, 1 player, imperfect information?
  - One, two or
  - Perfect info Deterministic, >1 player, perfect information? Deterministic, >1 player, imperfect information?
- Want algorit which recon Stochastic, 1 player, perfect information? Stochastic, 1 player, imperfect information?

Stochastic, >1 player, perfect information? Stochastic, >1 player, imperfect information?

http://u.hal3.name/ic.pl?q=game

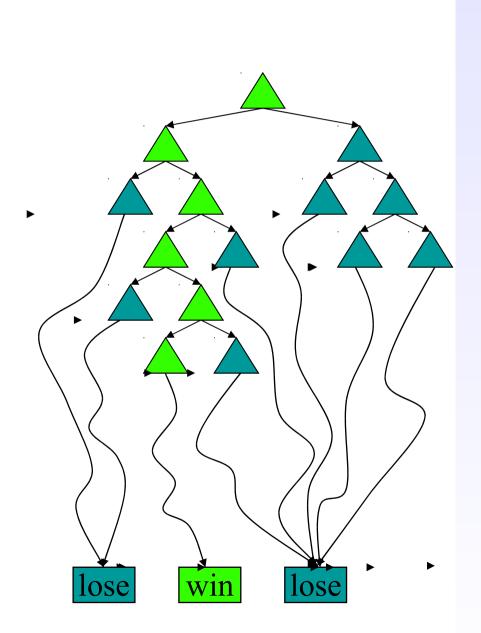
## **Deterministic Games**

- Many possible formalizations, one is:
  - > States: S (start at  $s_0$ )
  - Players: P={1...N} (usually take turns)
  - Actions: A (may depend on player / state)
  - ► Transition Function:  $SxA \rightarrow S$
  - ▶ Terminal Test:  $S \rightarrow \{t, f\}$
  - ► Terminal Utilities:  $SxP \rightarrow R$

### ► Solution for a player is a policy: $S \rightarrow A$

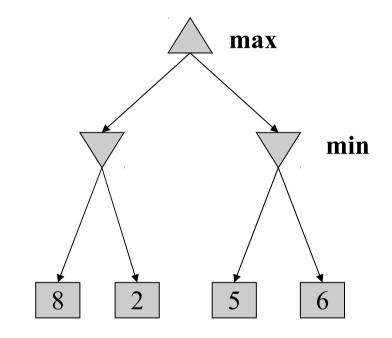
# **Deterministic Single-Player?**

- Deterministic, single player, perfect information:
  - Know the rules
  - Know what actions do
  - Know when you win
  - E.g. Freecell, 8-Puzzle, Rubik's cube
- … it's just search!
- Slight reinterpretation:
   Each node stores a value:
  - Each node stores a value: the best outcome it can reach
  - This is the maximal outcome of its children
  - Note that we don't have path sums as before (utilities at end)
- After search, can pick move that leads to best node

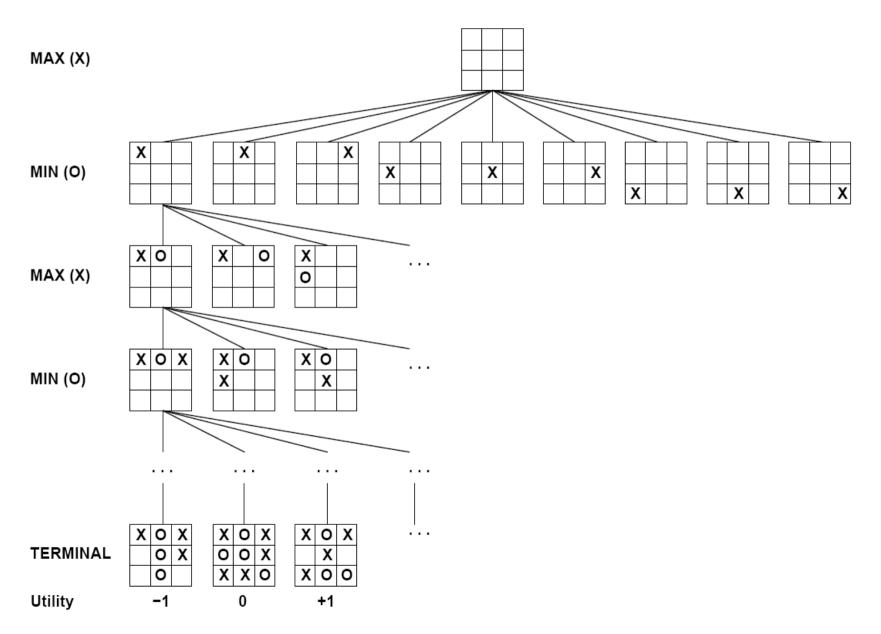


# **Deterministic Two-Player**

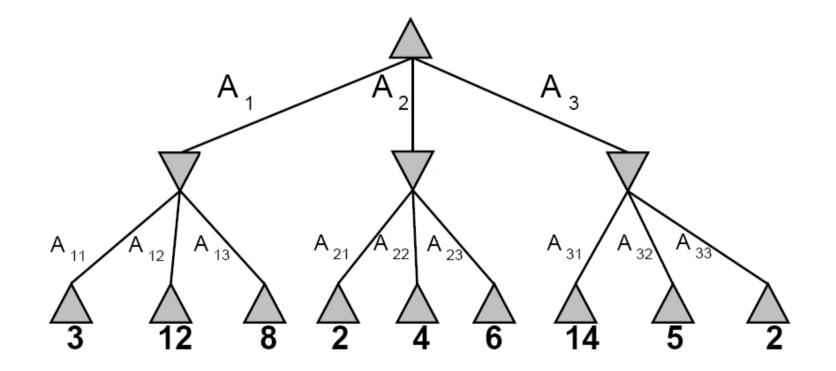
- E.g. tic-tac-toe, chess, checkers
- Minimax search
  - A state-space search tree
  - Players alternate
  - Each layer, or ply, consists of a round of moves
  - Choose move to position with highest minimax value = best achievable utility against best play
- Zero-sum games
  - One player maximizes result
  - The other minimizes result



### **Tic-tac-toe Game Tree**



## **Minimax Example**



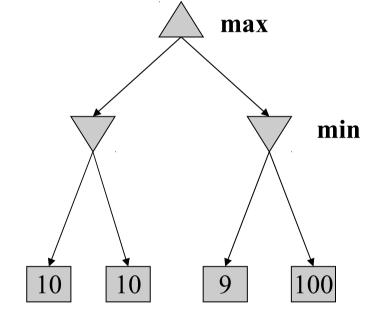
function MAX-VALUE(state) returns a utility value if TERMINAL-TEST(state) then return UTILITY(state)  $v \leftarrow -\infty$ for a, s in SUCCESSORS(state) do  $v \leftarrow MAX(v, MIN-VALUE(s))$ return v

function MIN-VALUE(state) returns a utility value if TERMINAL-TEST(state) then return UTILITY(state)  $v \leftarrow \infty$ for a, s in SUCCESSORS(state) do  $v \leftarrow MIN(v, MAX-VALUE(s))$ return v

## **Minimax Properties**

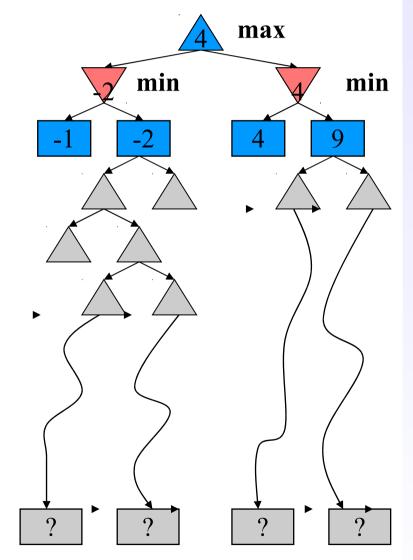
Optimal against a perfect player. Otherwise?

- Time complexity?
   O(b<sup>m</sup>)
- Space complexity?
  - O(bm)
- For chess,  $b \approx 35$ ,  $m \approx 100$ 
  - Exact solution is completely infeasible
  - But, do we need to explore the whole tree?



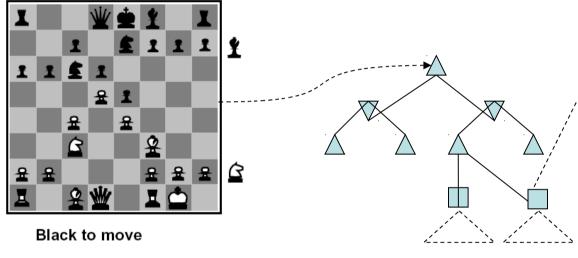
# **Resource Limits**

- Cannot search to leaves
- Depth-limited search
  - Instead, search a limited depth of tree
  - Replace terminal utilities with an eval function for non-terminal positions
- Guarantee of optimal play is gone
- More plies makes a BIG difference
  - [DEMO: limitedDepth]
- Example:
  - Suppose we have 100 seconds, can explore 10K nodes / sec
  - So can check 1M nodes per move
  - α-β reaches about depth 8 decent chess program

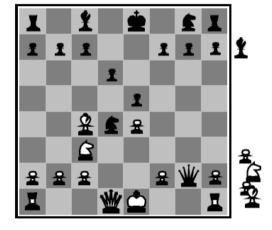


# **Evaluation Functions**

#### Function which scores non-terminals







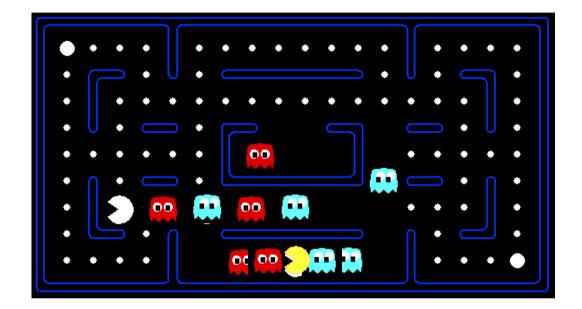
White to move Black winning

- Ideal function: returns the utility of the position
- In practice: typically weighted linear sum of features:

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \ldots + w_n f_n(s)$$

> e.g.  $f_1(s)$  = (num white queens – num black queens), etc.

### **Evaluation for Pacman**



[DEMO: thrashing, smart ghosts]

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \ldots + w_n f_n(s)$$

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# **Iterative Deepening**

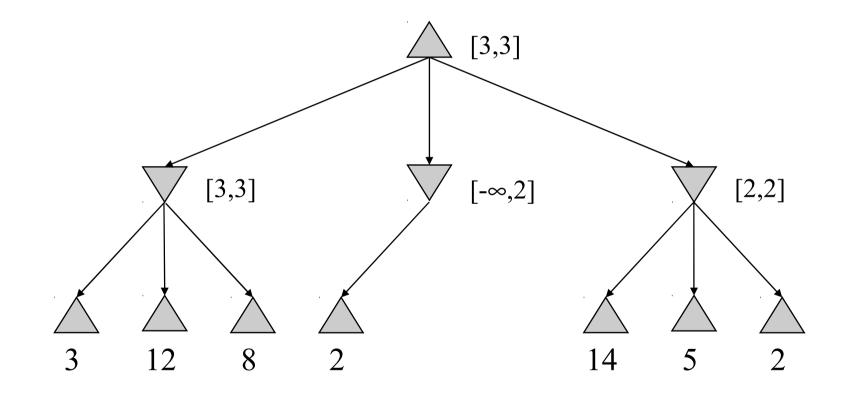
Iterative deepening uses DFS as a subroutine:

- Do a DFS which only searches for paths of length 1 or less. (DFS gives up on any path of length 2)
- 2. If "1" failed, do a DFS which only searches paths of length 2 or less.
- 3. If "2" failed, do a DFS which only searches paths of length 3 or less. ....and so on.

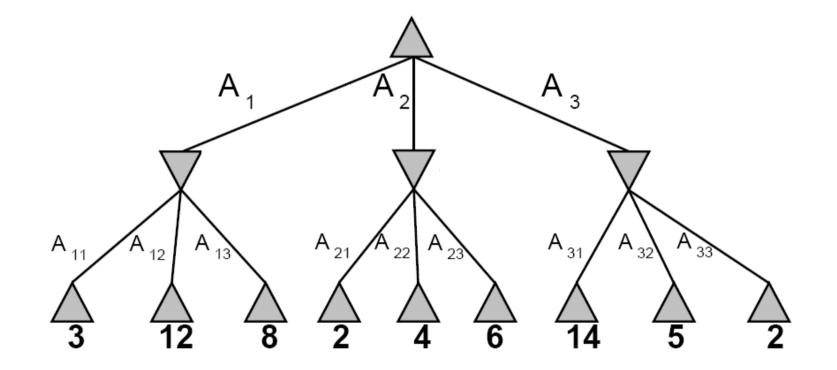
This works for single-agent search as well! Why do we want to do this for multiplayer games?

b	

## **Pruning in Minimax Search**



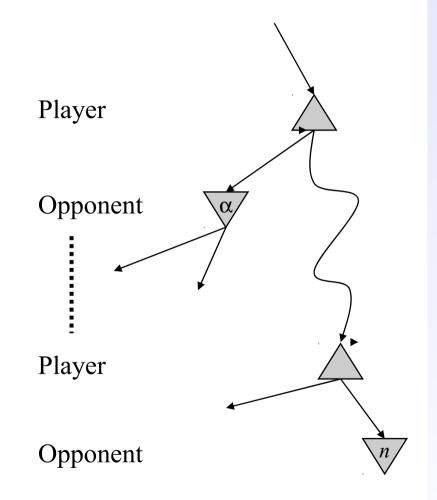
# $\alpha$ - $\beta$ Pruning Example



# $\alpha$ - $\beta$ Pruning

#### General configuration

- α is the best value that MAX can get at any choice point along the current path
- If *n* becomes worse than α,
   MAX will avoid it, so can stop considering *n*'s other children
- > Define  $\beta$  similarly for MIN



# α-β Pruning Pseudocode

function MAX-VALUE(state) returns a utility value
if TERMINAL-TEST(state) then return UTILITY(state)

 $v \leftarrow -\infty$ 

for a, s in SUCCESSORS(state) do  $v \leftarrow Max(v, MIN-VALUE(s))$ return v

function MAX-VALUE(state,  $\alpha$ ,  $\beta$ ) returns a utility value

**inputs**: *state*, current state in game

 $\alpha$ , the value of the best alternative for MAX along the path to state

 $\beta,$  the value of the best alternative for  $_{\rm MIN}$  along the path to state

if TERMINAL-TEST(*state*) then return UTILITY(*state*)

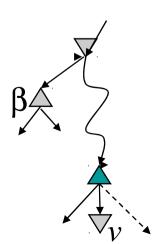
 $v \leftarrow -\infty$ 

for a, s in SUCCESSORS(state) do

 $v \leftarrow MAX(v, MIN-VALUE(s, \alpha, \beta))$ if  $v > \beta$  then return v

at 
$$v \geq \beta$$
 then return  $i \alpha \leftarrow MAX(\alpha, v)$ 

return v

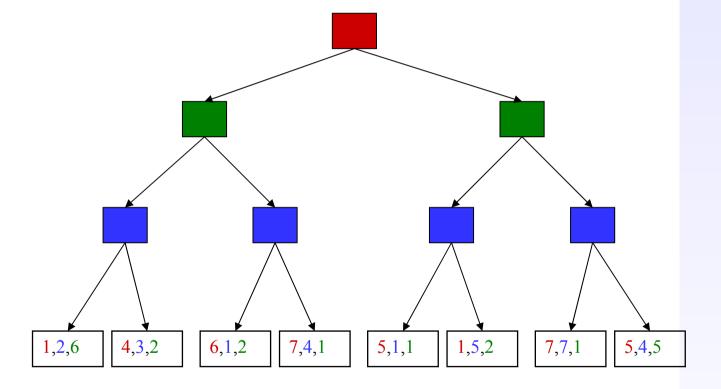


# **α-β Pruning Properties**

- Pruning has no effect on final result
- Good move ordering improves effectiveness of pruning
- With "perfect ordering":
  - Time complexity drops to O(b<sup>m/2</sup>)
  - Doubles solvable depth
  - > Full search of, e.g. chess, is still hopeless!
- A simple example of metareasoning, here reasoning about which computations are relevant

## **Non-Zero-Sum Games**

- Similar to minimax:
  - Utilities are now tuples
  - Each player maximizes their own entry at each node
  - Propagate (or back up) nodes from children

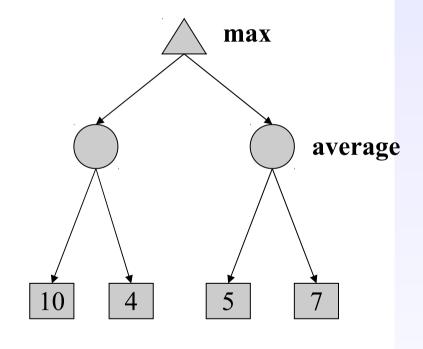


# **Stochastic Single-Player**

- What if we don't know what the result of an action will be? E.g.,
  - In solitaire, shuffle is unknown
  - In minesweeper, mine locations
  - In pacman, ghosts!

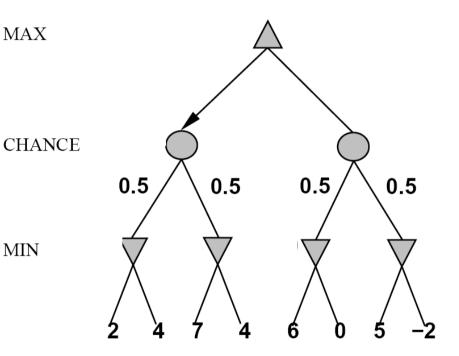
#### Can do expectimax search

- Chance nodes, like actions except the environment controls the action chosen
- Calculate utility for each node
- Max nodes as in search
- Chance nodes take average (expectation) of value of children
- Later, we'll learn how to formalize this as a Markov Decision Process



# **Stochastic Two-Player**

- E.g. backgammon
- Expectiminimax (!)
  - Environment is an extra player that moves after each agent
  - Chance nodes take expectations, otherwise like minimax



 $\mathbf{if}\ state\ \mathbf{is}\ \mathbf{a}\ \mathrm{MAX}\ \mathbf{node}\ \mathbf{then}$ 

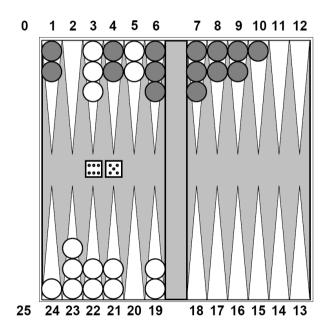
return the highest EXPECTIMINIMAX-VALUE of SUCCESSORS(*state*) if *state* is a MIN node then

**return** the lowest EXPECTIMINIMAX-VALUE of SUCCESSORS(*state*) **if** *state* is a chance node **then** 

 ${f return}$  average of  ${f Expect}MINIMAX$ -VALUE of  ${f Successors}({\it state})$ 

## **Stochastic Two-Player**

- Dice rolls increase b: 21 possible rolls with 2 dice
  - Backgammon ≈ 20 legal moves
  - Depth 4 = 20 x (21 x 20)<sup>3</sup> 1.2 x 10<sup>9</sup>
- As depth increases, probability of reaching a given node shrinks
  - So value of lookahead is diminished
  - So limiting depth is less damaging
  - But pruning is less possible...
- TDGammon uses depth-2 search + very good eval function + reinforcement learning: worldchampion level play



# What's Next?

### Make sure you know what:

- Probabilities are
- Expectations are
- You should be able to do any exercise from:
  - http://www.cs.umd.edu/class/fall2011/cmsc250-0x0x/hw/HW11.pdf
  - Username and password are both "250"
- If you can't, review your probability discrete math!
  - http://www.cs.umd.edu/class/fall2011/cmsc250-0x0x/notes/CRASH.pdf

### > Next topics:

- Dealing with uncertainty
- How to learn evaluation functions
- Markov Decision Processes